Applications of the Channel Theory

Commonsense Reasoning Representations

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## Overview

- Repetition
- Commonsense Reasoning / Nonmonotonicity
- (Imperfect) Representations


## Basics

- tokens (particulars, instances): things in the world (in time) - $a, b, c$
- types (in state spaces: states): $a, \beta$, ?, $s$
- classification: $\boldsymbol{A}$, set of tokens $(A)$ is classified in set of types
- if a is of type $a$ we write: $a_{1}{ }_{A} a$ (with respect to $A$ )


## Information Channels and Local Logics

- C is an information channel, consists of core $\boldsymbol{C}$ and infomorphisms from parts to $\boldsymbol{C}$ :

- L is a local logic, consists of $\boldsymbol{A}$, a set of constraints of $\boldsymbol{A}$ and a set of normal tokens
- normal tokens: satisfy all constraints in L


## State Spaces

- $\boldsymbol{S}$ is a state space, consists of a set of tokens $(S)$, a set of states $\left(\mathrm{O}_{S}\right)$ and a function mapping between them: $\boldsymbol{S}=\left\langle S, \mathrm{O}_{S}\right.$, state $\rangle$
- $\operatorname{Evt}(\boldsymbol{S})$ is the according classification
- $\log (\boldsymbol{S})$ is the local logic on $\operatorname{Evt}(\boldsymbol{S})$


## Overview

- Problem of Nonmonotonicity
- State Spaces, enhanced
- Background Conditions
- Relativising to a Background Condition

Baseball:
$\left(a_{I}\right)$ pitcher throws ball to batter $(\beta)$ ball will arrive at batter

$$
\Rightarrow a_{1}+\beta
$$

monotonicity:
$a_{1}, a_{2}+\beta$
but:
$\left(a_{2}\right)$ ball hits bird

$$
\Rightarrow a_{1}, a_{2}+\neg \beta
$$



## Real valued State Spaces

- $\boldsymbol{S}=\left\langle S, \mathrm{O} \subseteq \mathrm{R}^{n}\right.$, state $\rangle$
- each state is a vector $s$ with dimension $n$
- $s$ has input and output coordinates ( $s=$ $s_{i} \times s_{o}$ ) $=$ observables
- outputs can be computed from inputs


## Judith's heating system

- a state $s=\left(s_{1}, s_{2}, s_{3}, s_{4}, s_{5}, s_{6}, s_{7}\right)$



## Background Conditions

- $B$ is a function from domain $P$ to real numbers
- $P$ in inputs of $S$
- is called set of parameters of $B$
- a state $s$ satisfies $B$ if the corresponding inputs of $s$ have same value as given by $B$ ( $s_{i}=B(i) \forall i \in P$ )
- $B_{1}=B_{2} \Leftrightarrow P_{1} \subseteq P_{2}$


## Judith's heating system

( $a_{I}$ ) thermostat: $65=s_{I}=70$

$$
a_{l}, a_{2}+\beta
$$

$\left(a_{2}\right)$ room temperature: $s_{2}=58$

$$
a_{1}, a_{2}, a_{3}+\neg \beta
$$

$\left(a_{3}\right)$ power: $s_{3}=0$
$(\beta)$ hot air is coming out of the vents
$\Rightarrow a_{1}, a_{2}, \beta$ are silent about $s_{3}, s_{4}, s_{5}$ (which are supposed to be the parameters)
but:
$a_{3}$ is not silent about $s_{3}$
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## Silence

- $a$ is silent on $B$, if it does not tell anything about the parameters of B :
$-s=_{B} s^{\prime}$ if $s_{t}=s_{t}{ }^{\prime} \forall t \notin B$
$-\forall s, s^{\prime}$ : if $s={ }_{B} s^{\prime}$ and $s \in a$ then $s^{‘} \in a$
$\Rightarrow$ if we are reasoning about an observable $t$, then $t$ must be either an explicit input or output of the system (and not a parameter)
$\qquad$


## Judith's heating system

$\left(a_{1}\right)$ thermostat: $65=s_{1}=70$
$\left(a_{2}\right)$ room temperature: $s_{2}=58$
$\left(a_{3}\right)$ power: $s_{3}=0$
$(\beta)$ hot air is coming out of the vents
$\mathrm{B}=\left\{s_{3}=s_{4}=s_{5}=1\right\}$
$\Rightarrow B ? a_{3}=\left\{s_{4}=s_{5}=1\right\}$

## Weakening

- Gis a set of types
- the weakening of $B$ by $G(B$ ? $G)$ is the greatest lower bound of all $B_{0}=B$ such that every type $a \in G$ is silent on $B_{0}$
$\Rightarrow B ? a$ is restriction of $B$ to the set of inputs $i \in P$ such that $a$ is silent on $i$


## Relativising to a Background Condition

- $S_{B}$ is relativisation of $S$ to $B$
- subspace of $\boldsymbol{S}$
- only states that satisfy $B$
- $\log \left(S_{B}\right)$ is the local logic on $\operatorname{Evt}(S)$ supported by $B$
- consistent states are those satisfying $B$
- entailment only over states satisfying B
- $G+_{B}$ ? $\Leftrightarrow \forall s$ sat. $B$ (if $s \in p \forall p \in G$ then $s \in q$ for some $q \in$ ?)
- normal tokens are those satisfying $B$


## Judith's heating system

$\left(a_{l}\right)$ thermostat: $65=s_{l}=70$
$\left(a_{2}\right)$ room temperature: $s_{2}=58$
$\left(a_{3}\right)$ power: $s_{3}=0$
$(\beta)$ hot air is coming out of the vents

- $a 1, a 2+\beta$ holds in $\log \left(S_{B}\right)$
- because $a_{3}$ not silent about $s_{3}$ : switch to $B$ ? $a_{3}$ $\Rightarrow a 1, a 2+\beta$ does not hold in $\log \left(S_{B ? a 3}\right)$ (is no constraint there)
$\Rightarrow a 1, a 2, a 3+\neg \beta$ is constraint in $\log \left(\boldsymbol{S}_{B ? a 3}\right)$


## Strict Entailment

- $G \Rightarrow_{B}$ ?: $G$ strictly entails ? relative to B
$-G t_{\log \left(S_{B}\right)}$ ?
- all types in $G \cup$ ? are silent on $B$
- then (conclusions):
$-\underset{?}{G} \Rightarrow_{B}$ ? is a better model of human reasoning than $G+_{B}$
$-G \Rightarrow_{B}$ ? is monotonic in $G$ and ? (only with weakening)
- if you have a type $a$ not silent on $B$ it is natural to weaken $B$ : $B$ ? a
$-G \Rightarrow_{B}$ ? does not entail: $G, a \Rightarrow_{B ? a}$ ? or $G \Rightarrow_{B \backslash a}$ ?, $a$


## Representations

## Overview

- The Problem of Imperfect Representations
- Representation Systems
- Explaining Imperfect Representations



## The bridge




## Representation Systems

- $R=\langle C, L\rangle$ is a representation system
$-\mathrm{C}=\{f: \boldsymbol{A} \leftrightarrows \boldsymbol{C}, \mathrm{g}:$ $\boldsymbol{B} \leftrightarrows C\}$ is a binary channel
$-L$ is the local logic on the core $\boldsymbol{C}$


Representations Targets

## Representations

- $a$ is a representation of $b$, if $a, b$ are connected by some $c \in \boldsymbol{C}$
- $a$ is accurate representation of $b$, if $c$ is normal token $\left(c \in N_{\mathrm{L}}\right)$
- content of $a: a_{1}^{1} G$ $f[G]+\llcorner g(?)$
- $a$ represents $b$ as being of type $\beta$, if $\beta$ in content of $a$
if $a$ is accurate representation of $b$ and $a$ represents $b$ as being of type $\beta$, then $b_{\left.\right|_{B}} \beta$


## Explaining Imperfect <br> Representations

- tokens in target classification ( $\boldsymbol{B}$ ) are really regions at times
- if $b_{0}$ changes to $b_{1}$ this gives rise to a new connection $c_{1}$ between $a$ and $b_{1}$
$\Rightarrow a$ represents both: $b_{0}$ and $b_{1}$
$\Rightarrow$ but $c_{1}$ supports not all of the constraints of $R$



## References

- Jon Barwise and Jerry Seligman, Information Flow, The Logic of Distributed Systems, Cambridge University Press, 1997

